贵州大学 计算机科学与技术学院 实验报告

| 院(系)名称 | 示范性软件学院 | 班级 | 软工 206 | 课程名称 | Linux 系统 |
|--------|----------------|------------|--------|-------|----------|
| 实验名称 | 实验四 | 日期 | 9月18日 | 指导教师 | 王老师 |
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1. 实验名称 进程管理

2. 实验目的

熟悉 Linux 进程创建系统调用 fork、vfork、clone;认识进程的地址空间(共享的以及独占的);进一步认识进程并发以及简单的进程调试。

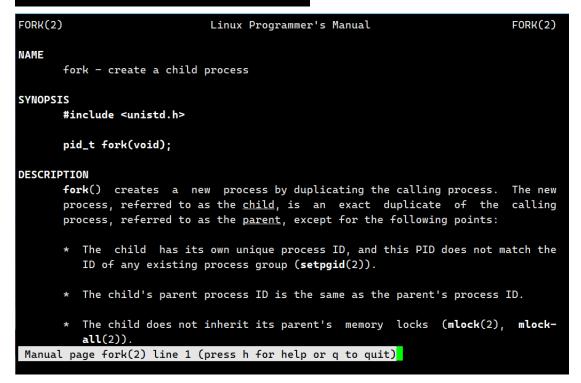
2. 内容(原理、方法)

分别用 fork、vfork、clone 创建进程;查看父子和兄弟进程中变量值的变化情况,查看进程中变量的逻辑地址和物理地址;查看进程的状态等信息。

3. 结果、分析与建议

- a) 了解 fork、vfork、clone 系统调用的区别;了解 gdb 的调试命令。
 - man fork

[root@hadoop100 ~]# man fork



- man vfork

[root@hadoop100 ~]# man vfork

```
VFORK(2)
                                                                          VFORK(2)
                             Linux Programmer's Manual
NAME
       vfork - create a child process and block parent
SYNOPSIS
       #include <sys/types.h>
       #include <unistd.h>
       pid_t vfork(void);
   Feature Test Macro Requirements for glibc (see feature_test_macros(7)):
       vfork():
           Since glibc 2.12:
              _BSD_SOURCE ||
                   (_XOPEN_SOURCE >= 500 ||
                       _XOPEN_SOURCE && _XOPEN_SOURCE_EXTENDED) &&
                   !(_POSIX_C_SOURCE >= 200809L || _XOPEN_SOURCE >= 700)
           Before glibc 2.12:
               _BSD_SOURCE || _XOPEN_SOURCE >= 500 ||
               _XOPEN_SOURCE && _XOPEN_SOURCE_EXTENDED
Manual page vfork(2) line 1 (press h for help or q to quit)
```

- man clone

[root@hadoop100 ~]# man clone

```
CLONE(2)
                                                                          CLONE(2)
                            Linux Programmer's Manual
NAME
      clone, __clone2 - create a child process
SYNOPSIS
       /* Prototype for the glibc wrapper function */
      #define _GNU_SOURCE
      #include <sched.h>
      int clone(int (*fn)(void *), void *child_stack,
                int flags, void *arg, ...
                 /* pid_t *ptid, void *newtls, pid_t *ctid */ );
       /* For the prototype of the raw system call, see NOTES */
DESCRIPTION
      clone() creates a new process, in a manner similar to fork(2).
      This page describes both the glibc clone() wrapper function and the under-
      lying system call on which it is based. The main text describes the wrap-
Manual page clone(2) line 1 (press h for help or q to quit)
```

b) 实验步骤

i. 进程中变量的值和地址

代码:

```
ntptr_t mem_addr(unsigned long vaddr, unsigned long *paddr)
    int pagesize = getpagesize();
    unsigned long v_pageindex = vaddr / pagesize;
unsigned long v_offset = v_pageindex * sizeof(uint64_t);
    unsigned long page_offset = vaddr % pagesize;
    uint64_t item = 0;
int fd = open("/proc/self/pagemap", 0_RDONLY);
    lseek(fd, v_offset, SEEK_SET);
read(fd, &item, sizeof(uint64_t));
    if((((uint64_t)1 << 63) & item) == 0)</pre>
         printf("page present is 0\n");
return 0 ;
     uint64_t phy_pageindex = (((uint64_t)1 << 55) - 1) & item;</pre>
    *paddr = (phy_pageindex * pagesize) + page_offset;
    return *paddr;
    main(void)
    char str[10];
    int count = 1;
unsigned long pa[2]={0,0};
                          st.txt", O_RDWR);
     int fd = open('
     if(fork() == 0)
         printf("pid:%d, ppid:%d\n",getpid(), getppid());
read(fd, str, 10);
         count += 1;
printf("Child process
                                    : %s\n", (char *)str);
         mem_addr((unsigned long)&count, &pa[1]);
nrintf("virtual addr of str=%n and count
                                                                       ical addr of str=%p, &count=%p\n",str,&count, p
         printf("
a[0], pa[1]);
         printf("count: %d (%p), pid: %d\n", count, &count, getpid());
    else
         printf("pid:%d, ppid:%d\n",getpid(), getppid());
read(fd, str, 10);
                                                                                         str=%p, &count=%p\n",str,&count,
         printf('
mem_addr((intptr_t)str, &pa[0]), mem_addr((intptr_t)&count, &pa[1]));
                     Father process : %s\n", (char *)str);
count: %d (%p), pid: %d\n", count, &count, getpid());
         printf("
         printf(
    sleep(10);
    return 0;
"proc-1.c" 67L, 1843C
                                                                                                          18,5
                                                                                                                          顶端
```

编译运行: 普通用户运行

```
[q@hadoop100 experiment_004]$ ./proc-1
pid:3941, ppid:3890
virtual addr of str=0x7ffe5e93d2d0 and &count=0x7ffe5e93d2cc, physical addr of str=0x2d0, &count=0x2cc
Father process :
@
count: 1 (0x7ffe5e93d2cc), pid: 3941
pid:3942, ppid:3941
Child process :
@
virtual addr of str=0x7ffe5e93d2d0 and count=0x7ffe5e93d2cc, physical addr of str=0x2d0, &count=0x2cc
count: 2 (0x7ffe5e93d2cc), pid: 3942
```

root 超级用户运行

```
[root@hadoop100 experiment_004]# gcc -g proc-1.c -o proc-1
[root@hadoop100 experiment_004]# ./proc-1
pid:3843, ppid:2605
virtual addr of str=0x7ffda00492b0 and &count=0x7ffda00492ac, physical addr of str=0x5c9562b0, &count=0x
5c9562ac
Father process :
@
count: 1 (0x7ffda00492ac), pid: 3843
pid:3844, ppid:3843
Child process :
@
virtual addr of str=0x7ffda00492b0 and count=0x7ffda00492ac, physical addr of str=0x8e75c2b0, &count=0x8
e75c2ac
count: 2 (0x7ffda00492ac), pid: 3844
```

ii. 将上面程序 fork 替换为 vfork,并以普通用户和超级用户运行,检查输出的结果。

代码:

```
tptr_t mem_addr(unsigned long vaddr, unsigned long *paddr)
    int pagesize = getpagesize();
    unsigned long v_pageindex = vaddr / pagesize;
unsigned long v_offset = v_pageindex * sizeof(uint64_t);
     unsigned long page_offset = vaddr % pagesize;
    uint64_t item = 0;
int fd = open("/proc/self/pagemap", 0_RDONLY);
    lseek(fd, v_offset, SEEK_SET);
read(fd, &item, sizeof(uint64_t));
     if((((uint64_t)1 << 63) & item) == 0)</pre>
         printf("page present is 0\n");
return 0 ;
     uint64_t phy_pageindex = (((uint64_t)1 << 55) - 1) & item;</pre>
     *paddr = (phy_pageindex * pagesize) + page_offset;
     return *paddr;
    main(void)
     char str[10];
    int count = 1;
unsigned long pa[2]={0,0};
     int fd = open("test.txt", O_RDWR);
if(vfork() == 0)
         printf("pid:%d, ppid:%d\n",getpid(), getppid());
read(fd, str, 10);
         count += 1;
printf("Child process : %s\n", (char *)str);
mem_addr((unsigned long)str, &pa[0]);
         mem_addr((unsigned long)&count, &pa[1]);
                                                               sp, physical addr of str=%p, &count=%p\n",str,&count, p
         printf("
a[0], pa[1]);
         printf("count: %d (%p), pid: %d\n", count, &count, getpid());
     else
         printf("pid:%d, ppid:%d\n",getpid(), getppid());
read(fd, str, 10);
                                                                                              %p, &count=%p\n",str,&count,
         printf("
mem_addr((intptr_t)str, &pa[0]), mem_addr((intptr_t)&count, &pa[1]));
                    Father process : %s\n", (char *)str);
count: %d (%p), pid: %d\n", count, &count, getpid());
         printf("
          printf('
    sleep(10);
   插入 ---
                                                                                                            67,1
                                                                                                                            底端
```

普通用户运行

```
[root@hadoop100 experiment_004]# gcc -g proc-2.c -o proc-2
[root@hadoop100 experiment_004]# su q
[q@hadoop100 experiment_004]$ ./proc-2
pid:4098, ppid:4097
Child process:
@
virtual addr of str=0x7ffdbf686370 and count=0x7ffdbf68636c, physical addr of str=0x370, &count=0x36c
count: 2 (0x7ffdbf68636c), pid: 4098

pid:4097, ppid:4060
virtual addr of str=0x7ffdbf686370 and &count=0x7ffdbf68636c, physical addr of str=0x370, &count=0x36c
Father process:

count: 0 (0x7ffdbf68636c), pid: 4097
pid:4105, ppid:4097
Child process:
virtual addr of str=0x7ffdbf686370 and count=0x7ffdbf68636c, physical addr of str=0x370, &count=0x36c
count: 2 (0x7ffdbf68636c), pid: 4105

ptitual addr of str=0x7ffdbf686370 and count=0x7ffdbf68636c, physical addr of str=0x370, &count=0x36c
count: 2 (0x7ffdbf68636c), pid: 4105

ptitual addr of str=0x7ffdbf68636c), pid: 4105
```

root 超级用户运行

iii. ps 如何工作

```
[root@hadoop100 experiment_004]# ps
PID TTY TIME CMD
2605 pts/0 00:00:00 bash
4200 pts/0 00:00:00 ps
[root@hadoop100 experiment_004]#
```

```
[root@hadoop100 experiment_004]# strace -e file -o x ps -l 2605
F S UID PID PPID C PRI NI ADDR SZ WCHAN TTY TIME CMD
4 S 0 2605 2596 0 80 0 - 29135 do_wai pts/0 0:00 -bash
[root@hadoop100 experiment_004]#
```

```
[root@hadoop100 experiment_004]# grep 2605 x
execve("/usr/bin/ps", ["ps", "-l", "2605"], 0x7ffeff417bd0 /* 26 vars */) = 0
stat("/proc/2605", {st_mode=S_IFDIR|0555, st_size=0, ...}) = 0
open("/proc/2605/stat", 0_RDONLY) = 6
open("/proc/2605/status", 0_RDONLY) = 6
open("/proc/2605/cmdline", 0_RDONLY) = 6
open("/proc/2605/wchan", 0_RDONLY) = 6
readlink("/proc/2605/fd/2", "/dev/pts/0", 127) = 10
[root@hadoop100 experiment_004]#
```

[root@hadoop100 experiment_004]# cat /proc/2605/stat
2605 (bash) S 2596 2605 2605 34816 4240 4202752 5144 119584 1 46 6 12 128 217 20 0 1 0 8408 119336960 80
6 18446744073709551615 4194304 5100836 140722994558496 140722994557144 140613722527244 0 65536 3686404 1
266761467 18446744072460308870 0 0 17 3 0 0 1 0 0 7200240 7236240 39456768 140722994560641 1407229945606
47 140722994560647 140722994561006 0
[root@hadoop100 experiment_004]#

[root@hadoop100 experiment_004]# man proc

| proc(3tcl) | | Tcl Built-In | Commands | | proc(3tcl) |
|------------------|----------------------|-----------------|---------------|------------|--------------------|
| | | | | | |
| NAME | | | | | |
| proc - 建고 | ☑一个 Tcl 过程 | | | | |
| 总览 SYNOPSIS | | | | | |
| proc <u>name</u> | <u>args body</u> | | | | |
| 描述 DESCRIPTION | | | | | |
| | 令建立一个叫做name | 的新的 Tcl | 过程, 恭捧已经 | | 何现存的命令或过程. |
| | | | | | 是未限定的(unquali- |
| | 括任何包含(这个过程) | | | | |
| | 字空间限定符(qualifi | | | | // XE/IE VANC Hame |
| | | | | | 参数。每个参数指定符(|
| spec- | 2)III)//F(2) % - G | , , 0-04011/047 | 31,0±, 410 | 41/02/17/2 | 2 M. 4 / 2 M.H.C. |
| i- | | | | | |
| fier)也可以 | J是有一个或两个字段(· | field)的一个列表 | 。如果在指定往 | 符中只有一个单一 | 字段则它是参数的名字 |
| ; 如果有两个字段, | 则第一个是参数名而第 | 二个是它的缺省值 | 直。 | | |
| | | | | | |
| 在调用 | | | | | <u>name</u> |
| 时,为过程 | 的每个形式参数建立一 | 个局部变量;它的 | 值将是在调用包 | 命令中相应的(实图 | 示)参数的值或这个参数 |
| 的缺省值。在过程说 | 間用中可以不指定有缺省 | î值的参数。但是, | 必须有足够的 | 实际参数给所有? | 没有缺省值的形式参数, |
| 并且没有多余的实际 | 示参数。有一种特殊情况 | 2可以允许过程有可 | 可变数目的参数 | [。如果最后的形式 | 式参数的名字是 |
| args,则到 | 这个过程的一个调用包 | 含的实际参数可以 | 多于过程拥有的 | 的形式参数。此时 | ,把开始于应当被赋给 |
| Manual page proc | (3tcl) line 1/59 54% | (press h for he | lp or q to qu | uit) | |

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```
[root@hadoop100 experiment_004]# cat /proc/2605/status
Name:
       bash
Umask:
       0022
State:
       S (sleeping)
Tgid:
       2605
Ngid:
       2605
Pid:
PPid:
       2596
TracerPid:
              0
Uid:
       0
             0
                    0
                           0
Gid:
       0
              0
FDSize: 256
Groups: 0
VmPeak:
        116540 kB
VmSize:
        116540 kB
VmLck:
            0 kB
VmPin:
             0 kB
VmHWM:
          3224 kB
VmRSS:
          3224 kB
RssAnon:
                 1372 kB
RssFile:
                 1852 kB
RssShmem:
                    0 kB
VmData:
          1284 kB
VmStk:
           132 kB
VmExe:
           888 kB
VmLib:
          2148 kB
VmPTE:
            60 kB
             0 kB
VmSwap:
Threads:
SigQ: 0/15603
SigPnd: 00000000000000000
ShdPnd: 00000000000000000
SigBlk: 000000000010000
SigIgn: 000000000384004
SigCgt: 000000004b813efb
CapInh: 0000000000000000
CapPrm: 0000001fffffffff
CapEff: 0000001fffffffff
CapBnd: 0000001fffffffff
CapAmb: 0000000000000000
NoNewPrivs:
Seccomp:
Speculation_Store_Bypass:
                           thread vulnerable
Mems_allowed_list:
voluntary_ctxt_switches:
nonvoluntary_ctxt_switches:
[root@hadoop100 experiment_004]#
```

4. 附录(如源程序)

源代码以上截屏均有